

Introduction to Computers and Programming

Prof. I. K. Lundqvist

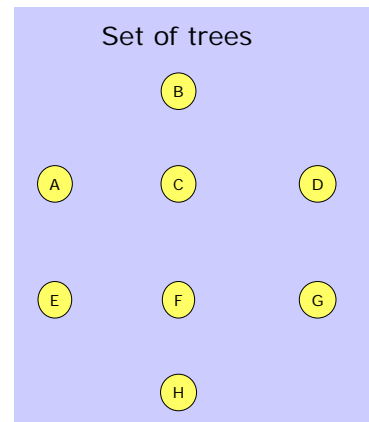
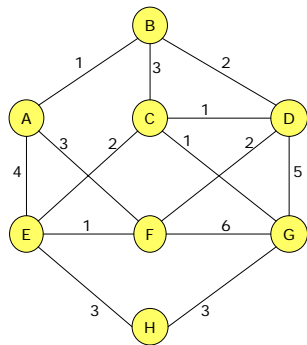
Recitation 2
April 6 2004

Minimum Spanning Tree

Kruskal's Algorithm

– Finds a minimum spanning tree for a connected weighted graph

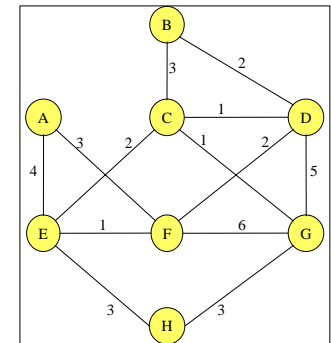
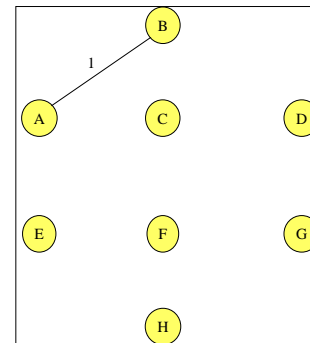
- Create a set of trees, where each vertex in the graph is a separate tree
- Create set S containing all edges in the graph
- While S not empty
 - Remove edge with minimum weight from S
 - if that edge connects two different trees, then add it to the forest, combining two trees into a single tree
 - Otherwise discard that edge



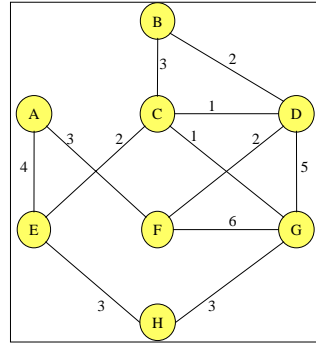
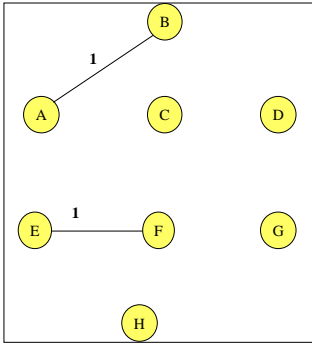
| S | A | B | C | D | E | F | G | H |
|---|---|---|---|---|---|---|---|---|
| A | 0 | | | | | | | |
| B | 1 | 0 | | | | | | |
| C | | 3 | 0 | | | | | |
| D | | 2 | 2 | 0 | | | | |
| E | 4 | | 2 | | 0 | | | |
| F | 3 | | | 2 | 1 | 0 | | |
| G | | | 1 | 5 | | 6 | 0 | |
| H | | | | | 3 | | 3 | 0 |

Kruskal's Algorithm

Step 1

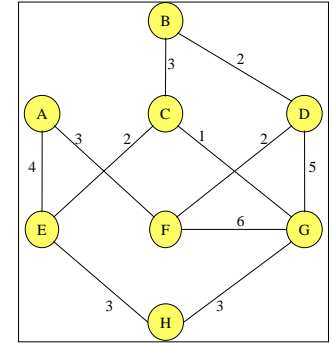
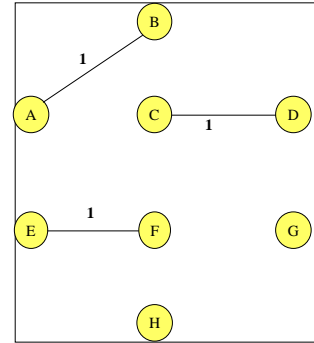


Step 2



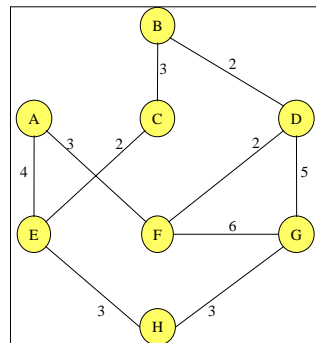
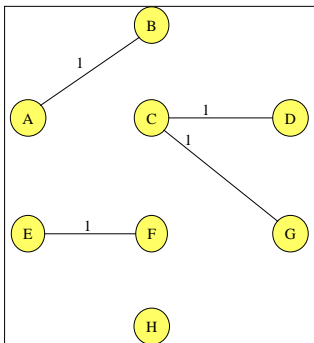
5

Step 3



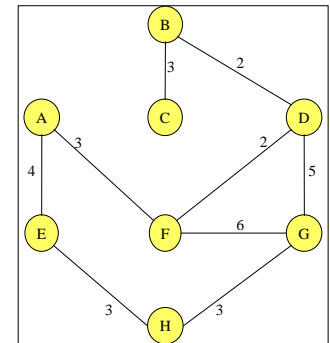
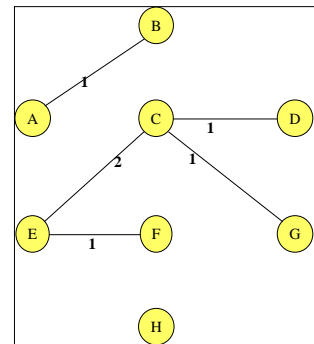
6

Step 4



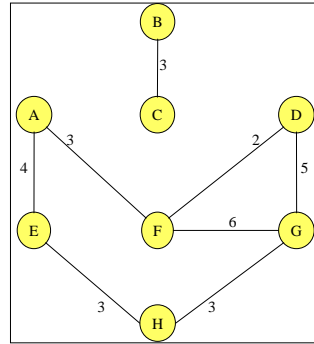
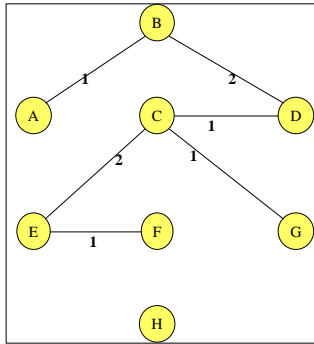
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Step 5



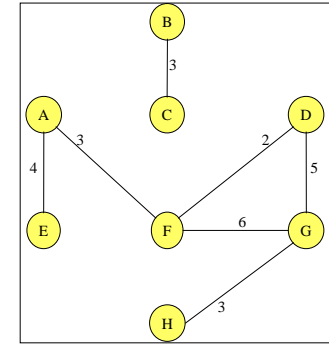
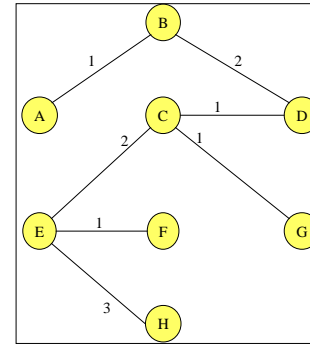
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Slide 6



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Slide 7



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Minimum Spanning Tree

• Prim's Algorithm

- Finds a subset of the edges (that form a tree) including every vertex and the total weight of all the edges in tree is minimized

- Choose starting vertex
- Create the Fringe Set

Initialization

- Loop until the MST contains all the vertices in the graph
 - Remove edge with minimum weight from Fringe Set
 - Add the edge to MST
 - Update the Fringe Set

Body

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Prim – Initialization

- Pick any vertex x as the starting vertex
- Place x in the Minimum Spanning Tree (MST)
- For each vertex y in the graph that is adjacent to x
 - Add y to the Fringe Set
- For each vertex y in the Fringe Set
 - Set weight of y to weight of the edge connecting y to x
 - Set x to be parent of y

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Prim – Body

While number of vertices in MST < vertices in the graph

Find vertex y with minimum weight in the Fringe Set

Add vertex and the edge x,y to the MST

Remove y from the Fringe Set

For all vertices z adjacent to y that are not in MST

If z is not in the Fringe Set

Add z to the Fringe Set

Set parent to y

Set weight of z to weight of the edge connecting z to y

Else

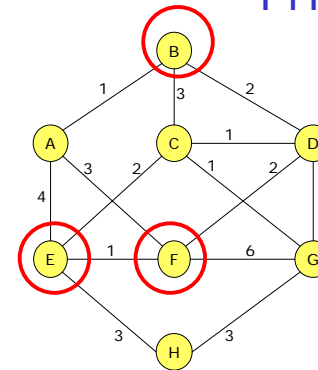
If $\text{Weight}(y,z) < \text{Weight}(z)$ then

Set parent to y

Set weight of z to weight of the edge connecting z to y

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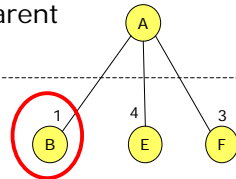
Prim's Algorithm



MST



Parent

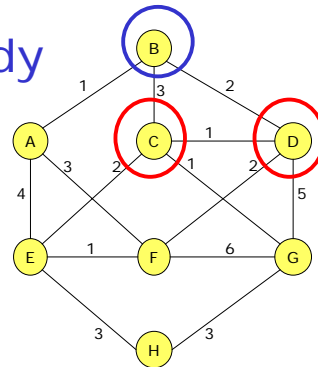
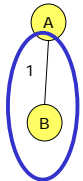


Fringe Set

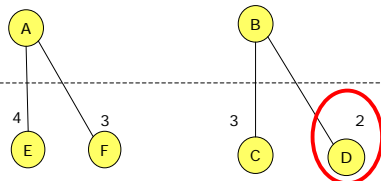
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Prim's Body

MST



Parent

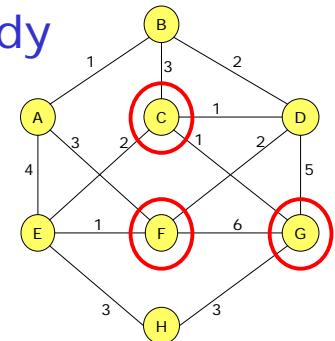
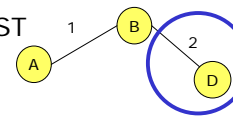


Fringe Set

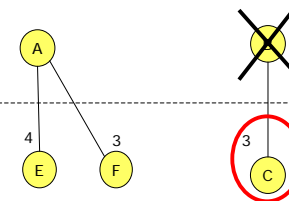
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Prim's Body

MST

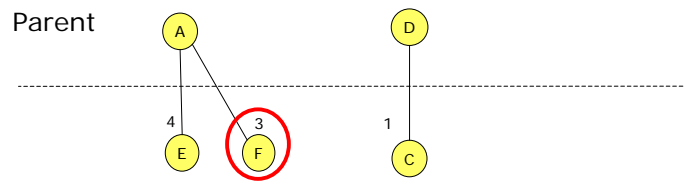
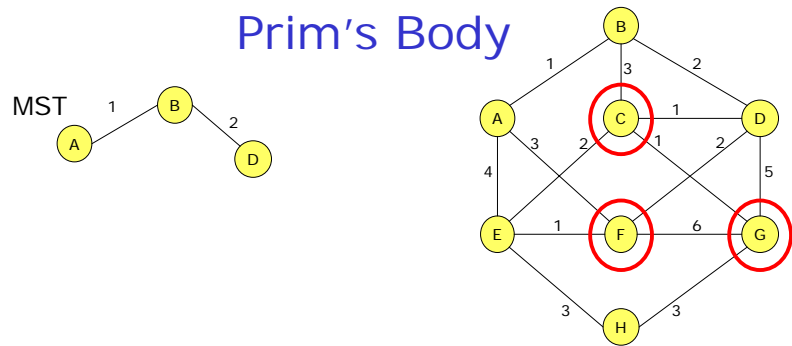


Parent



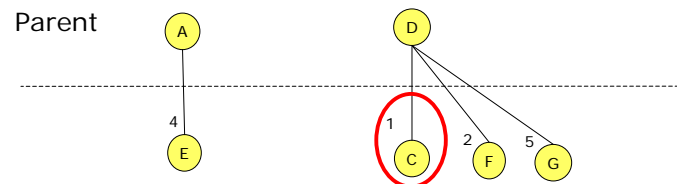
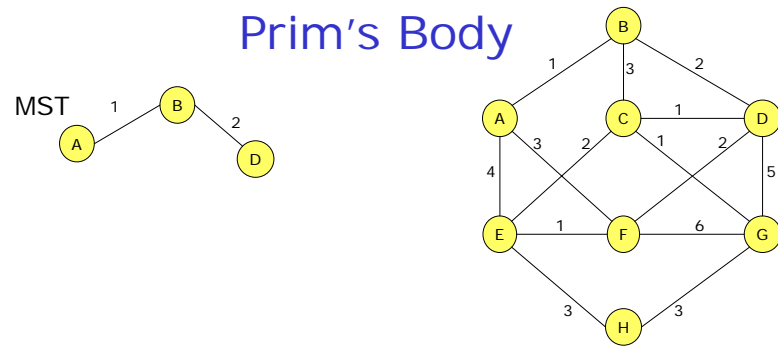
Fringe Set

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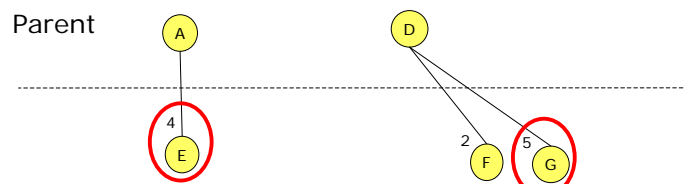
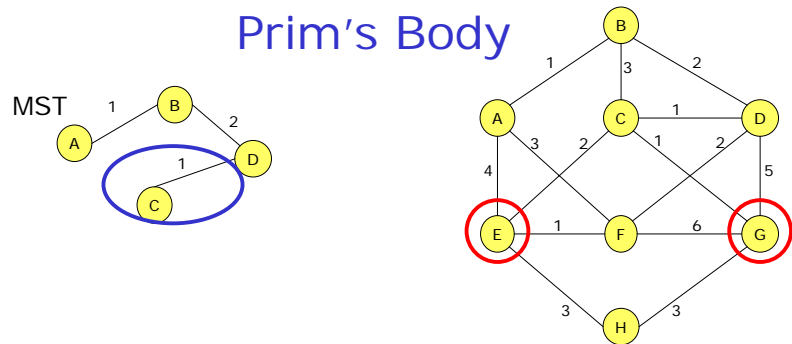
Fringe Set

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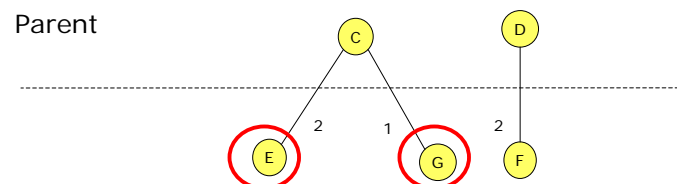
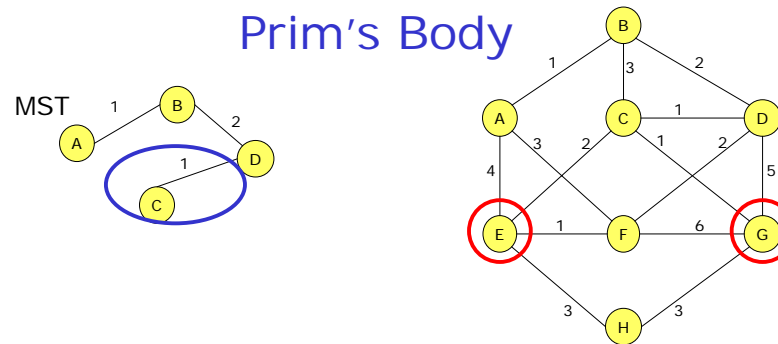
Fringe Set

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Fringe Set

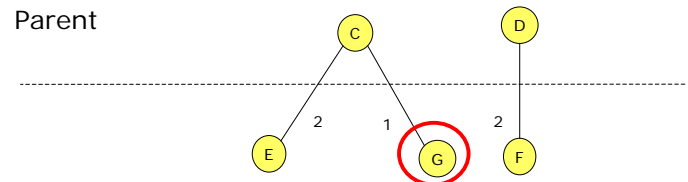
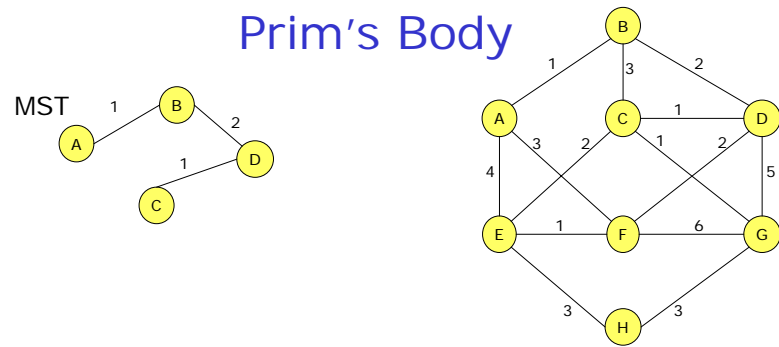
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Fringe Set

20

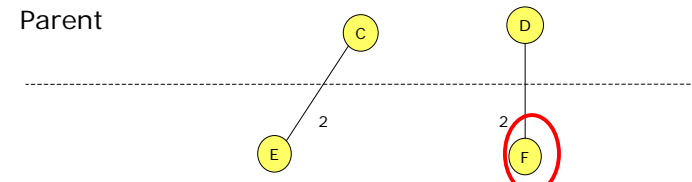
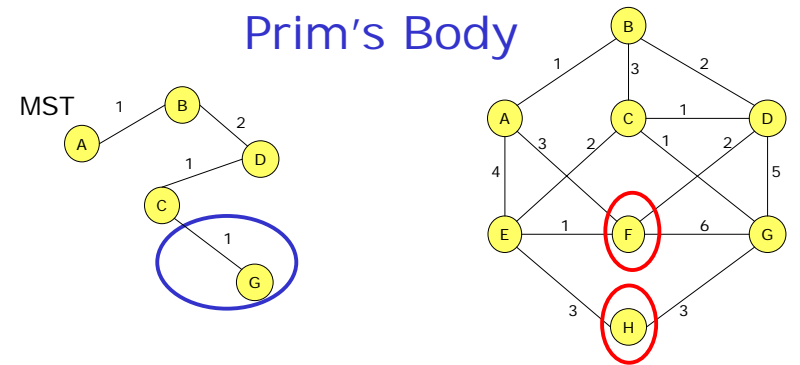
Prim's Body



Fringe Set

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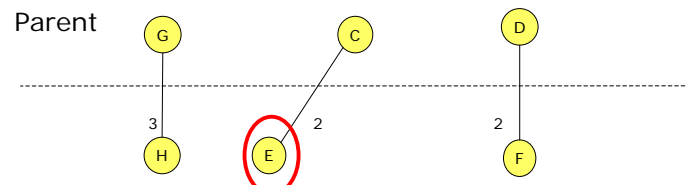
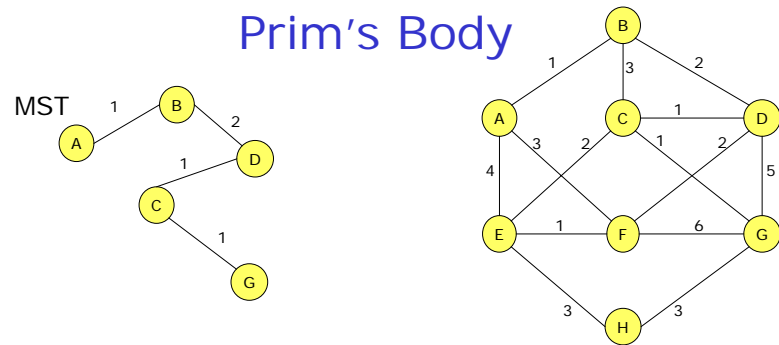
Prim's Body



Fringe Set

22

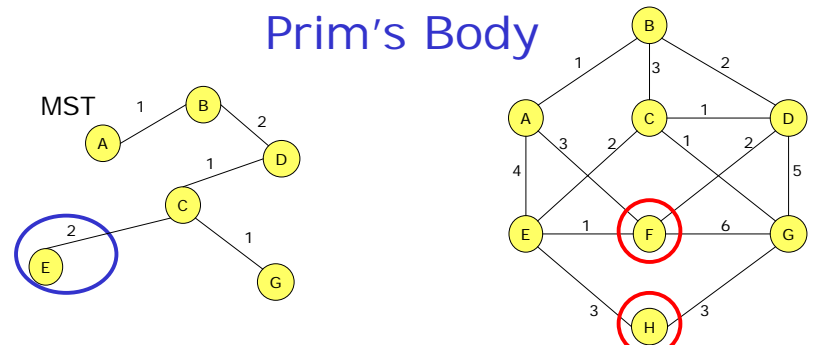
Prim's Body



Fringe Set

23

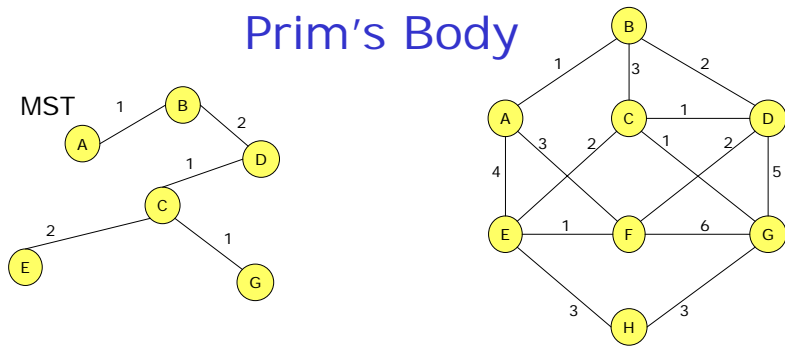
Prim's Body



Fringe Set

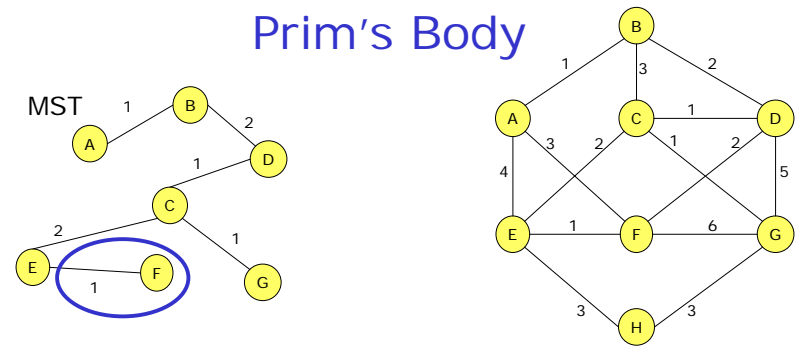
24

Prim's Body



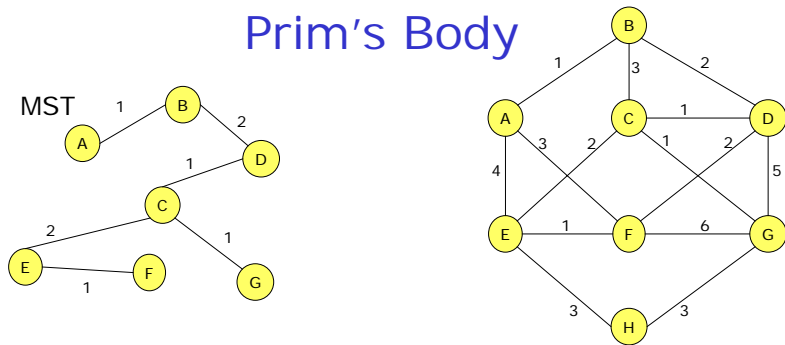
25

Prim's Body



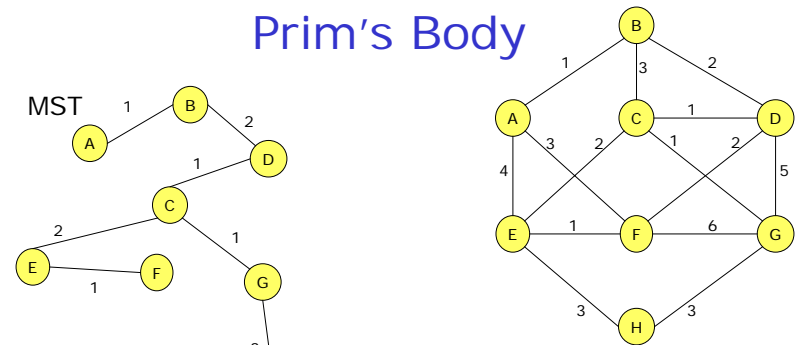
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Prim's Body



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Prim's Body



DONE!

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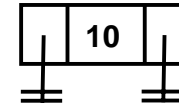
Inserting into Ordered Binary Tree

```
-- insert procedure
procedure Insert (Root      : in out Nodeptr;
                  Element   : in      Integer ) is
    New_Node : Nodeptr;
begin
    if Root = null then
        New_Node := new Node;
        New_Node.Element := Element;
        Root := New_Node;
    else
        if Root.Element < Element then
            Insert(Root.Right_Child, Element);
        else
            Insert(Root.Left_Child, Element);
        end if;
    end if;
end Insert;
```

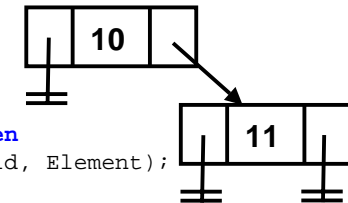
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Inserting into a Binary node

- Insert 10, 11, 9, 7, 8, 12
- Insert 10



- Insert 11

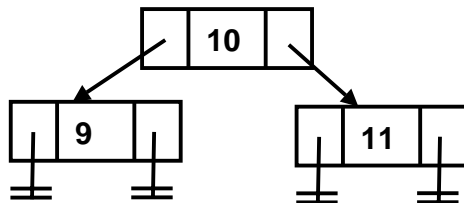


```
if Root.Element < Element then
    Insert(Root.Right_Child, Element);
else
    Insert(Root.Left_Child, Element);
end if;
```

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Inserting into Ordered Binary Tree

- Insert 9

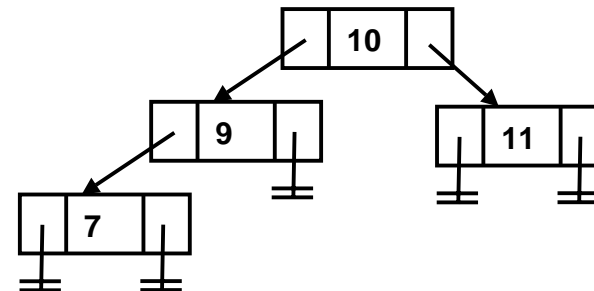


```
if Root.Element < Element then
    Insert(Root.Right_Child, Element);
else
    Insert(Root.Left_Child, Element);
end if;
```

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Inserting into Ordered Binary Tree

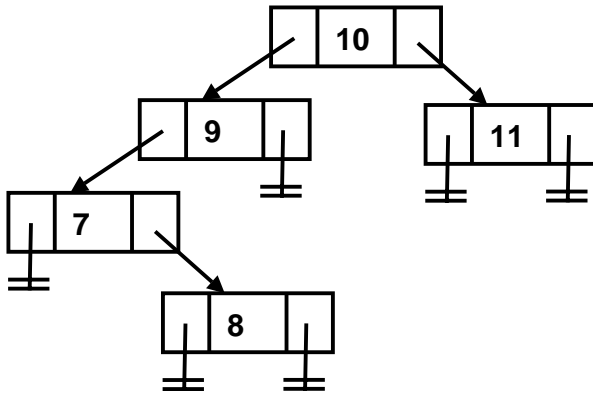
- Insert 7



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Inserting into Ordered Binary Tree

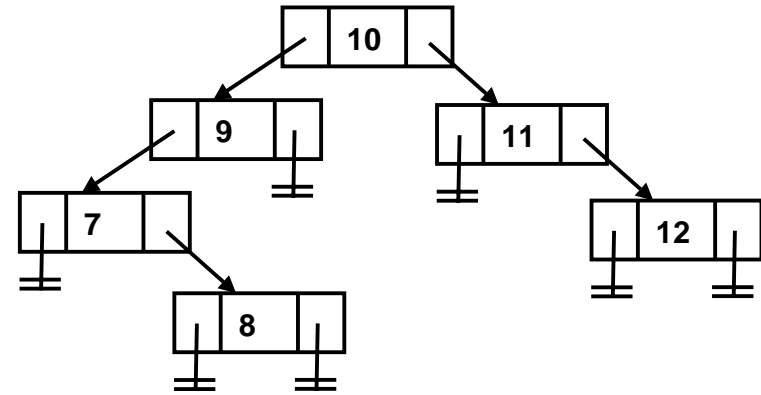
- Insert 8



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Inserting into Ordered Binary Tree

- Insert 12



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Shortest Path Problems

- **Dijkstra's algorithm**

- Finds shortest path for a directed and connected graph $G(V,E)$ which has non-negative weights.
- Applications:
 - Internet routing
 - Road generation within a geographic region
 - ...

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Dijkstra's Algorithm

- Dijkstra(G,w,s)

Init_Source(G,s)

$S :=$ empty set

$Q :=$ set of all vertices

while Q is not an empty set **loop**

$u :=$ Extract_Min(Q)

$S := S$ union $\{u\}$

for each vertex v which is a neighbor of u **loop**

 Relax(u,v,w)

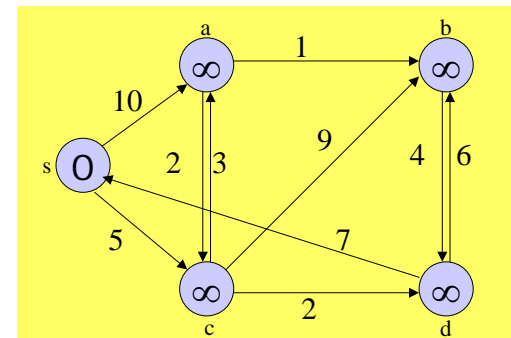
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Dijkstra's Algorithm

- **Init_Source(G,s)**
 for each vertex v in V[G] loop
 d[v] := infinite
 previous[v] := 0
 d[s] := 0
- v = **Extract_Min(Q)** searches for the vertex v in the vertex set Q that has the least d[v] value. That vertex is removed from the set Q and then returned.
- **Relax(u,v,w)**
 if d[v] > d[u] + w(u,v) then
 d[v] := d[u] + w(u,v)
 previous[v] := u

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Dijkstra's Algorithm



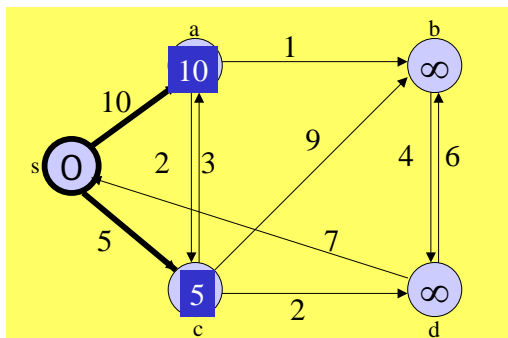
V = {a, b, c, d, s}
 E = {(s,c), (c,d), (d,b), (b,d), (c,b), (a,c), (c,a), (a,b), (s,a)}

S = {∅}
 Q = {s, a, b, c, d}

d = $\begin{pmatrix} 0 \\ \infty \\ \infty \\ \infty \\ \infty \end{pmatrix}$ prev = $\begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}$

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Dijkstra's Algorithm



S = {s}
 Q = {a, b, c, d}

d = $\begin{pmatrix} 0 \\ \infty \\ \infty \\ \infty \end{pmatrix} \rightarrow \begin{pmatrix} 0 \\ 10 \\ 5 \\ \infty \end{pmatrix}$

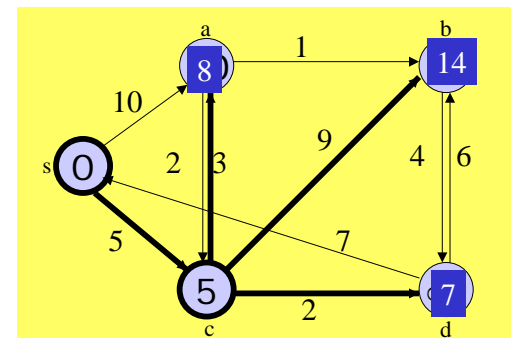
prev = $\begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \end{pmatrix} \rightarrow \begin{pmatrix} 0 \\ s \\ s \\ 0 \end{pmatrix}$

Extract_Min (Q) → s
 Neighbors of s = a, c

Relax (s,c,5)
 Relax (s,a,10)

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Dijkstra's Algorithm



S = {s, c}
 Q = {a, b, d}

d = $\begin{pmatrix} 0 \\ \infty \\ \infty \\ \infty \end{pmatrix} \rightarrow \begin{pmatrix} 0 \\ 8 \\ 14 \\ 5 \\ 7 \end{pmatrix}$

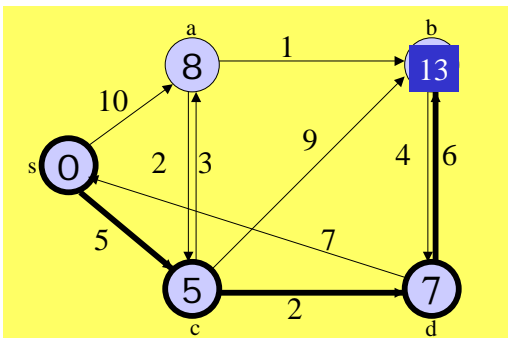
prev = $\begin{pmatrix} 0 \\ s \\ 0 \\ s \\ 0 \end{pmatrix} \rightarrow \begin{pmatrix} 0 \\ c \\ c \\ s \\ c \end{pmatrix}$

Extract_Min (Q) → c
 Neighbors of c = a, b, d

Relax (c,a,3)
 Relax (c,b,9)
 Relax (c,d,2)

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Dijkstra's Algorithm



$S = \{s, c, d\}$
 $Q = \{a, b\}$

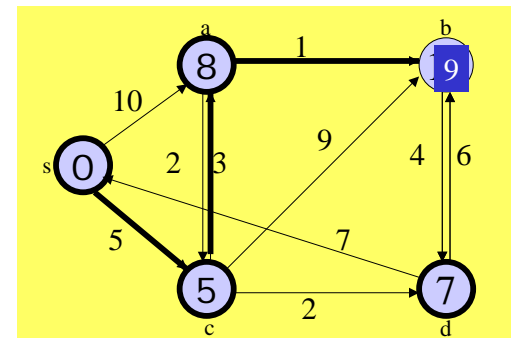
$d = \begin{pmatrix} 0 \\ 8 \\ 14 \\ 5 \\ 7 \end{pmatrix} \rightarrow \begin{pmatrix} 0 \\ 8 \\ 13 \\ 5 \\ 7 \end{pmatrix}$

$prev = \begin{pmatrix} 0 \\ c \\ c \\ s \\ c \end{pmatrix} \rightarrow \begin{pmatrix} 0 \\ c \\ d \\ s \\ c \end{pmatrix}$

Extract_Min (Q) \rightarrow d
 Neighbors of d = b
 Relax (d,b,6)

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Dijkstra's Algorithm



$S = \{s, c, d, a\}$
 $Q = \{b\}$

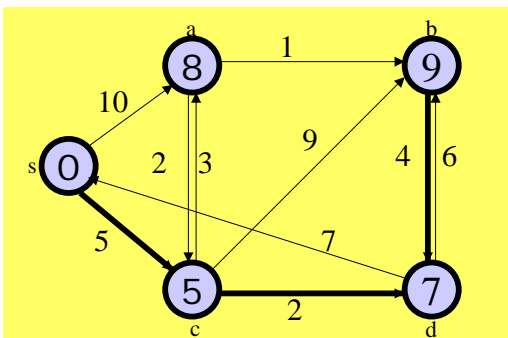
$d = \begin{pmatrix} 0 \\ 8 \\ 13 \\ 5 \\ 7 \end{pmatrix} \rightarrow \begin{pmatrix} 0 \\ 8 \\ 9 \\ 5 \\ 7 \end{pmatrix}$

$prev = \begin{pmatrix} 0 \\ c \\ d \\ s \\ c \end{pmatrix} \rightarrow \begin{pmatrix} 0 \\ c \\ a \\ s \\ c \end{pmatrix}$

Extract_Min (Q) \rightarrow a
 Neighbors of a = b, c
 Relax (a,b,1)
 Relax (a,c,3)

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Dijkstra's Algorithm



$S = \{s, c, d, a, b\}$
 $Q = \{\}$

$d = \begin{pmatrix} 0 \\ 8 \\ 9 \\ 5 \\ 7 \end{pmatrix} \rightarrow \begin{pmatrix} 0 \\ 8 \\ 9 \\ 5 \\ 7 \end{pmatrix}$

$prev = \begin{pmatrix} 0 \\ c \\ a \\ s \\ c \end{pmatrix} \rightarrow \begin{pmatrix} 0 \\ c \\ a \\ s \\ c \end{pmatrix}$

Extract_Min (Q) \rightarrow b
 Neighbors of b = d
 Relax (b, d, 4)

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