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16.323 Principles of Optimal Control  
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## 16.323 Lecture 16

### Model Predictive Control

- Allgower, F., and A. Zheng, Nonlinear Model Predictive Control, Springer-Verlag, 2000.
- Camacho, E., and C. Bordons, Model Predictive Control, Springer-Verlag, 1999.
- Kouvaritakis, B., and M. Cannon, Non-Linear Predictive Control: Theory & Practice, IEE Publishing, 2001.
- Maciejowski, J., Predictive Control with Constraints, Pearson Education POD, 2002.
- Rossiter, J. A., Model-Based Predictive Control: A Practical Approach, CRC Press, 2003.

- Planning in Lecture 8 was effectively “open-loop”
  - Designed the control input sequence  $\mathbf{u}(t)$  using an assumed model and set of constraints.
  - Issue is that with modeling error and/or disturbances, these inputs will not necessarily generate the desired system response.
- Need a “closed-loop” strategy to compensate for these errors.
  - Approach called **Model Predictive Control**
  - Also known as **receding horizon control**
- Basic strategy:
  - At time  $k$ , use knowledge of the system model to design an input sequence
 
$$\mathbf{u}(k|k), \mathbf{u}(k+1|k), \mathbf{u}(k+2|k), \mathbf{u}(k+3|k), \dots, \mathbf{u}(k+N|k)$$
 over a finite horizon  $N$  from the current state  $\mathbf{x}(k)$ 
    - Implement a fraction of that input sequence, usually just first step.
    - Repeat for time  $k+1$  at state  $\mathbf{x}(k+1)$

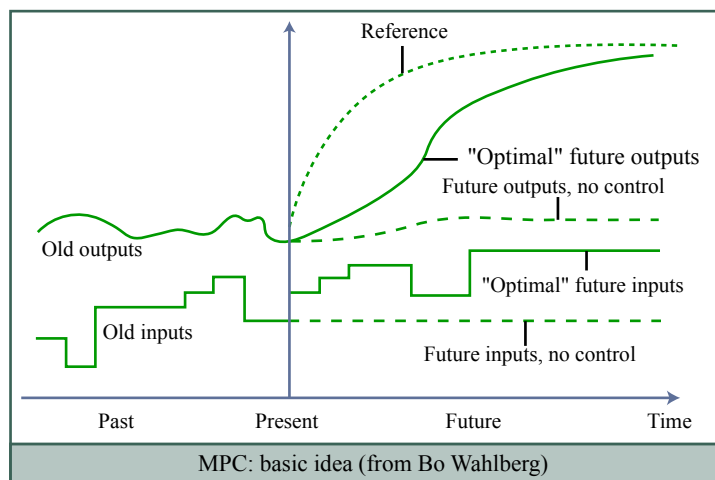


Figure by MIT OpenCourseWare.

- Note that the control algorithm is based on numerically solving an optimization problem at each step
  - Typically a constrained optimization
- Main advantage of MPC:
  - Explicitly accounts for system constraints.
    - ◇ Doesn't just design a controller to keep the system away from them.
  - Can easily handle nonlinear and time-varying plant dynamics, since the controller is explicitly a function of the model that can be modified in real-time (and plan time)
- Many commercial applications that date back to the early 1970's, see <http://www.che.utexas.edu/~qin/cpcv/cpcv14.html>
  - Much of this work was in process control - very nonlinear dynamics, but not particularly fast.
- As computer speed has increased, there has been renewed interest in applying this approach to applications with faster time-scale: trajectory design for aerospace systems.

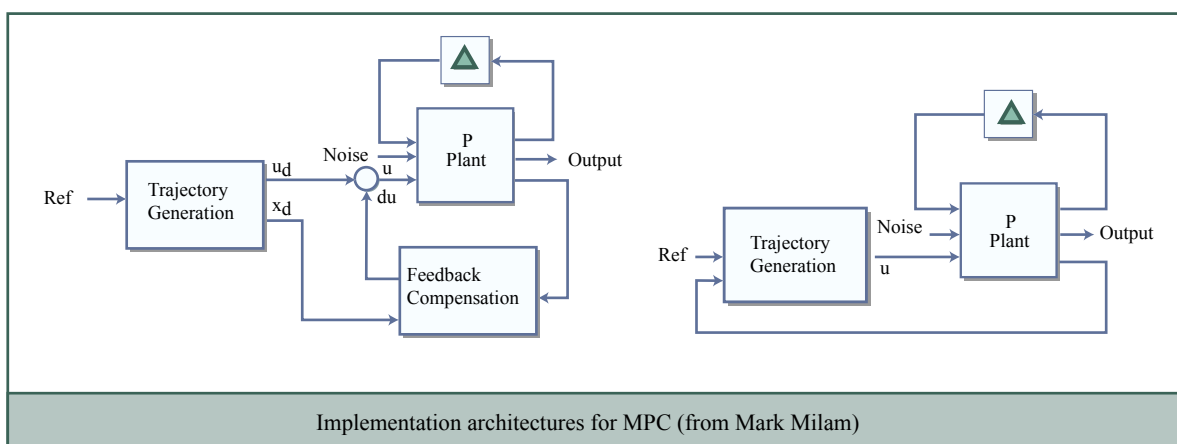


Figure by MIT OpenCourseWare.

- Given a set of plant dynamics (assume linear for now)

$$\begin{aligned}\mathbf{x}(k+1) &= A\mathbf{x}(k) + B\mathbf{u}(k) \\ \mathbf{z}(k) &= C\mathbf{x}(k)\end{aligned}$$

and a cost function

$$J = \sum_{j=0}^N \{ \|\mathbf{z}(k+j|k)\|_{R_{zz}} + \|\mathbf{u}(k+j|k)\|_{R_{uu}} \} + F(\mathbf{x}(k+N|k))$$

- $\|\mathbf{z}(k+j|k)\|_{R_{zz}}$  is just a short hand for a weighted norm of the state, and to be consistent with earlier work, would take

$$\|\mathbf{z}(k+j|k)\|_{R_{zz}} = \mathbf{z}(k+j|k)^T R_{zz} \mathbf{z}(k+j|k)$$

- $F(\mathbf{x}(k+N|k))$  is a terminal cost function

- Note that if  $N \rightarrow \infty$ , and there are no additional constraints on  $\mathbf{z}$  or  $\mathbf{u}$ , then this is just the discrete LQR problem solved on page 3–14.
  - Note that the original LQR result could have been written as just an input control sequence (**feedforward**), but we choose to write it as a linear state **feedback**.
  - In the nominal case, there is no difference between these two implementation approaches (feedforward and feedback)
  - But with modeling errors and disturbances, the state feedback form is much less sensitive.

⇒ This is the main reason for using feedback.

- Issue:** When limits on  $\mathbf{x}$  and  $\mathbf{u}$  are added, we can no longer find the general solution in analytic form ⇒ must solve it numerically.

- However, solving for a very long input sequence:
  - Does not make sense if one expects that the model is wrong and/or there are disturbances, because it is unlikely that the end of the plan will be implemented (a new one will be made by then)
  - Longer plans have more degrees of freedom and take much longer to compute.
  
- Typically design using a small  $N \Rightarrow$  short plan that does not necessarily achieve all of the goals.
  - Classical hard question is how large should  $N$  be?
  - If plan doesn't reach the goal, then must develop an estimate of the remaining **cost-to-go**
  
- Typical problem statement: for finite  $N$  ( $F = 0$ )

$$\min_u J = \sum_{j=0}^N \{ \|\mathbf{z}(k+j|k)\|_{R_{zz}} + \|\mathbf{u}(k+j|k)\|_{R_{uu}} \}$$

$$\text{s.t. } \mathbf{x}(k+j+1|k) = A\mathbf{x}(k+j|k) + B\mathbf{u}(k+j|k)$$

$$\mathbf{x}(k|k) \equiv \mathbf{x}(k)$$

$$\mathbf{z}(k+j|k) = C\mathbf{x}(k+j|k)$$

$$\text{and } |\mathbf{u}(k+j|k)| \leq u_m$$

- Consider converting this into a more standard optimization problem.

$$\mathbf{z}(k|k) = C\mathbf{x}(k|k)$$

$$\begin{aligned} \mathbf{z}(k+1|k) &= C\mathbf{x}(k+1|k) = C(A\mathbf{x}(k|k) + B\mathbf{u}(k|k)) \\ &= CA\mathbf{x}(k|k) + CB\mathbf{u}(k|k) \end{aligned}$$

$$\begin{aligned} \mathbf{z}(k+2|k) &= C\mathbf{x}(k+2|k) \\ &= C(A\mathbf{x}(k+1|k) + B\mathbf{u}(k+1|k)) \\ &= CA(A\mathbf{x}(k|k) + B\mathbf{u}(k|k)) + CB\mathbf{u}(k+1|k) \\ &= CA^2\mathbf{x}(k|k) + CAB\mathbf{u}(k|k) + CB\mathbf{u}(k+1|k) \\ &\vdots \end{aligned}$$

$$\begin{aligned} \mathbf{z}(k+N|k) &= CA^N\mathbf{x}(k|k) + CA^{N-1}B\mathbf{u}(k|k) + \dots \\ &\quad + CB\mathbf{u}(k+(N-1)|k) \end{aligned}$$

- Combine these equations into the following:

$$\begin{aligned} &\begin{bmatrix} \mathbf{z}(k|k) \\ \mathbf{z}(k+1|k) \\ \mathbf{z}(k+2|k) \\ \vdots \\ \mathbf{z}(k+N|k) \end{bmatrix} = \begin{bmatrix} C \\ CA \\ CA^2 \\ \vdots \\ CA^N \end{bmatrix} \mathbf{x}(k|k) \\ &+ \begin{bmatrix} 0 & 0 & 0 & \dots & 0 \\ CB & 0 & 0 & & 0 \\ CAB & CB & 0 & & 0 \\ \vdots & & & & \\ CA^{N-1}B & CA^{N-2}B & CA^{N-3}B & \dots & CB \end{bmatrix} \begin{bmatrix} \mathbf{u}(k|k) \\ \mathbf{u}(k+1|k) \\ \vdots \\ \mathbf{u}(k+N-1|k) \end{bmatrix} \end{aligned}$$

- Now define

$$Z(k) \equiv \begin{bmatrix} \mathbf{z}(k|k) \\ \vdots \\ \mathbf{z}(k+N|k) \end{bmatrix} \quad U(k) \equiv \begin{bmatrix} \mathbf{u}(k|k) \\ \vdots \\ \mathbf{u}(k+N-1|k) \end{bmatrix}$$

then, with  $\mathbf{x}(k|k) = \mathbf{x}(k)$

$$Z(k) = G\mathbf{x}(k) + HU(k)$$

- Note that

$$\sum_{j=0}^N \mathbf{z}(k+j|k)^T R_{zz} \mathbf{z}(k+j|k) = Z(k)^T W_1 Z(k)$$

with an obvious definition of the weighting matrix  $W_1$

- Thus

$$\begin{aligned} Z(k)^T W_1 Z(k) + U(k)^T W_2 U(k) &= (G\mathbf{x}(k) + HU(k))^T W_1 (G\mathbf{x}(k) + HU(k)) + U(k)^T W_2 U(k) \\ &= \mathbf{x}(k)^T H_1 \mathbf{x}(k) + H_2^T U(k) + \frac{1}{2} U(k)^T H_3 U(k) \end{aligned}$$

where

$$H_1 = G^T W_1 G, \quad H_2 = 2(\mathbf{x}(k)^T G^T W_1 H), \quad H_3 = 2(H^T W_1 H + W_2)$$

- Then the MPC problem can be written as:

$$\begin{aligned} \min_{U(k)} \tilde{J} &= H_2^T U(k) + \frac{1}{2} U(k)^T H_3 U(k) \\ \text{s.t.} \quad &\begin{bmatrix} I_N \\ -I_N \end{bmatrix} U(k) \leq u_m \end{aligned}$$



- **Key point:** the MPC problem is now in the form of a standard **quadratic program** for which standard and efficient codes exist.

QUADPROG Quadratic programming. %

X=QUADPROG(H,f,A,b) attempts to solve the %  
quadratic programming problem:

min  $0.5*x'*H*x + f'*x$  subject to:  $A*x \leq b$   
x

X=QUADPROG(H,f,A,b,Aeq,beq) solves the problem %  
above while additionally satisfying the equality%  
constraints  $Aeq*x = beq$ .

- Several Matlab toolboxes exist for testing these ideas
  - MPC toolbox by Morari and Ricker – extensive analysis and design tools.
  - MPCtools<sup>32</sup> enables some MPC simulation and is free  
[www.control.lth.se/user/johan.akesson/mpctools/](http://www.control.lth.se/user/johan.akesson/mpctools/)

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<sup>32</sup>Johan Akesson: "MPCtools 1.0 - Reference Manual". Technical report ISRN LUTFD2/TFRT-7613-SE, Department of Automatic Control, Lund Institute of Technology, Sweden, January 2006.

- Current form assumes that full state is available - can hookup with an estimator
  
- Current form assumes that we can sense and apply corresponding control immediately
  - With most control systems, that is usually a reasonably safe assumption
  - Given that we must re-run the optimization, probably need to account for this computational delay - different form of the discrete model - see F&P (chapter 2)
  
- If the constraints are not active, then the solution to the QP is that

$$U(K) = -H_3^{-1}H_2$$

which can be written as:

$$\begin{aligned}u(k|k) &= - [ 1 \ 0 \ \dots \ 0 ] (H^T W_1 H + W_2)^{-1} H^T W_1 G \mathbf{x}(k) \\ &= -K \mathbf{x}(k)\end{aligned}$$

which is just a state feedback controller.

- Can apply this gain to the system and check the eigenvalues.

- What can we say about the stability of MPC when the constraints are active? <sup>33</sup>
  - Depends a lot on the terminal cost and the terminal constraints. <sup>34</sup>
- Classic result: <sup>35</sup> Consider a MPC algorithm for a linear system with constraints. Assume that there are terminal constraints:
  - $\mathbf{x}(k + N|k) = 0$  for predicted state  $\mathbf{x}$
  - $\mathbf{u}(k + N|k) = 0$  for computed future control  $\mathbf{u}$
 Then if the optimization problem is feasible at time  $k$ ,  $\mathbf{x} = 0$  is stable.

**Proof:** Can use the performance index  $J$  as a Lyapunov function.

- Assume there exists a feasible solution at time  $k$  and cost  $J_k$
- Can use that solution to develop a feasible candidate at time  $k + 1$ , by simply adding  $\mathbf{u}(k + N + 1) = 0$  and  $\mathbf{x}(k + N + 1) = 0$ .
- **Key point:** can estimate the candidate controller performance

$$\begin{aligned}\tilde{J}_{k+1} &= J_k - \{ \|\mathbf{z}(k|k)\|_{R_{zz}} + \|\mathbf{u}(k|k)\|_{R_{uu}} \} \\ &\leq J_k - \{ \|\mathbf{z}(k|k)\|_{R_{zz}} \}\end{aligned}$$

- This candidate is suboptimal for the MPC algorithm, hence  $J$  decreases even faster  $J_{k+1} \leq \tilde{J}_{k+1}$
- Which says that  $J$  decreases if the state cost is non-zero (observability assumptions)  $\Rightarrow$  but  $J$  is lower bounded by zero.

- Mayne *et al.* [2000] provides excellent review of other strategies for proving stability – different terminal cost and constraint sets

<sup>33</sup>“Tutorial: model predictive control technology,” Rawlings, J.B. American Control Conference, 1999. pp. 662-676

<sup>34</sup>Mayne, D.Q., J.B. Rawlings, C.V. Rao and P.O.M. Scokaert, “Constrained Model Predictive Control: Stability and Optimality,” *Automatica*, 36, 789-814 (2000).

<sup>35</sup>A. Bemporad, L. Chisci, E. Mosca: “On the stabilizing property of SIORHC”, *Automatica*, vol. 30, n. 12, pp. 2013-2015, 1994.

- Consider a system similar to the Quansar helicopter<sup>36</sup>

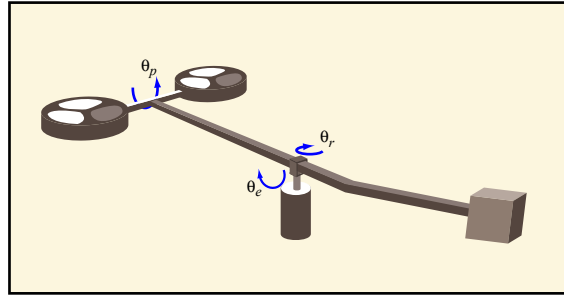


Figure by MIT OpenCourseWare.

- There are 2 control inputs – voltage to each fan  $V_f, V_b$
- A simple dynamics model is that:

$$\begin{aligned} \ddot{\theta}_e &= K_1(V_f + V_b) - T_g/J_e \\ \ddot{\theta}_r &= -K_2 \sin(\theta_p) \\ \ddot{\theta}_p &= K_3(V_f - V_b) \end{aligned}$$

and there are physical limits on the elevation and pitch:

$$-0.5 \leq \theta_e \leq 0.6 \quad -1 \leq \theta_p \leq 1$$

- Model can be linearized and then discretized  $T_s = 0.2\text{sec}$ .

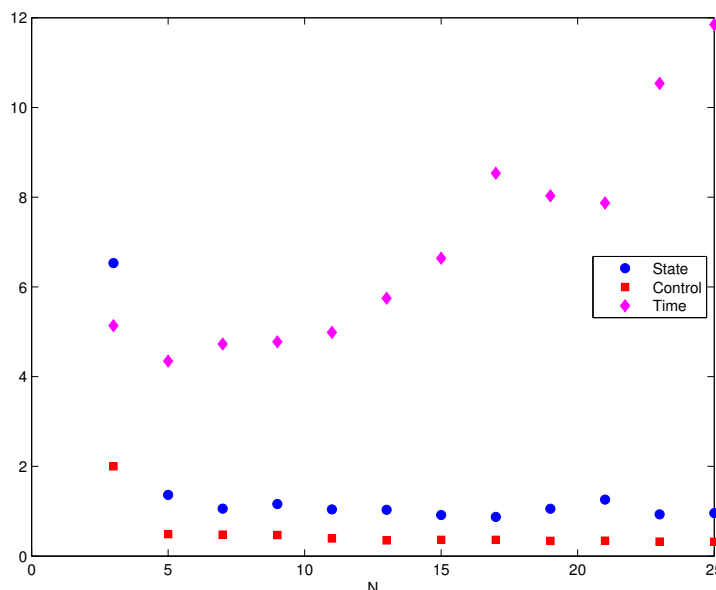


Figure 16.3: Response Summary

<sup>36</sup>ISSN 02805316 ISRN LUTFD2/TFRT- -7613- -SE MPCtools 1.0 Reference Manual Johan Akesson Department of Automatic Control Lund Institute of Technology January 2006

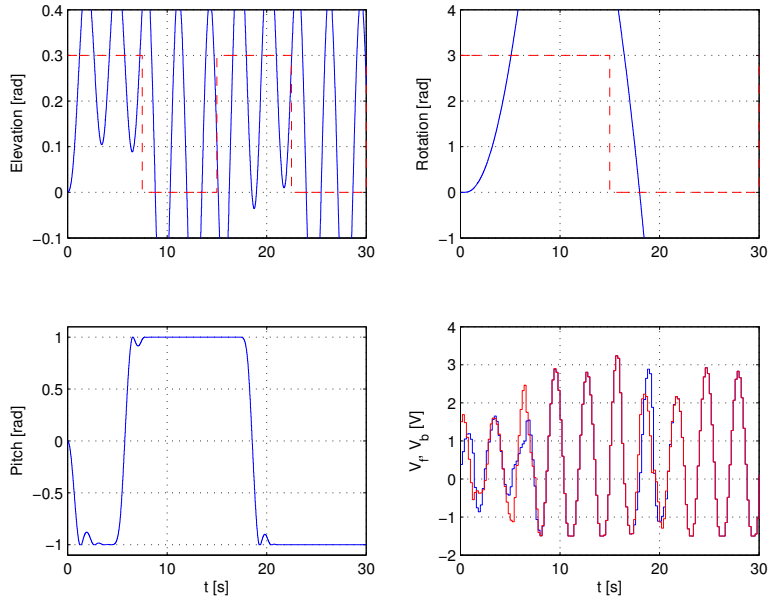


Figure 16.4: Response with  $N = 3$

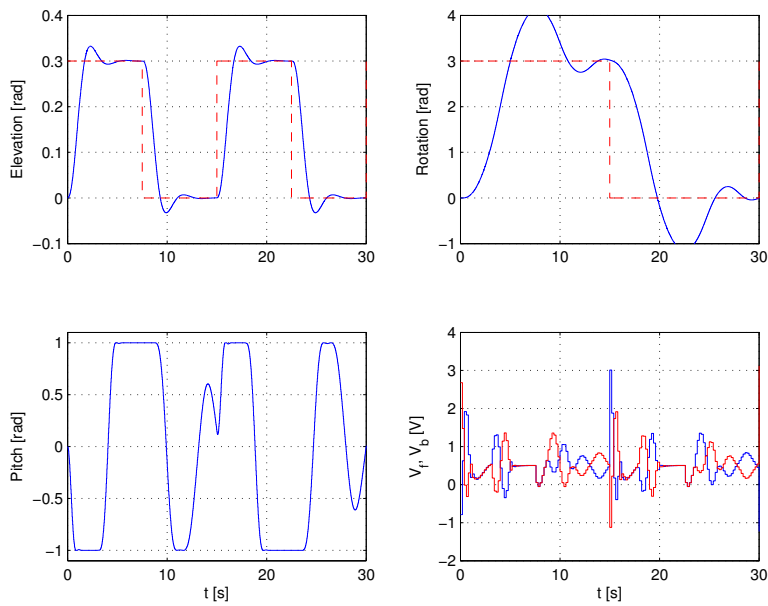


Figure 16.5: Response with  $N = 10$

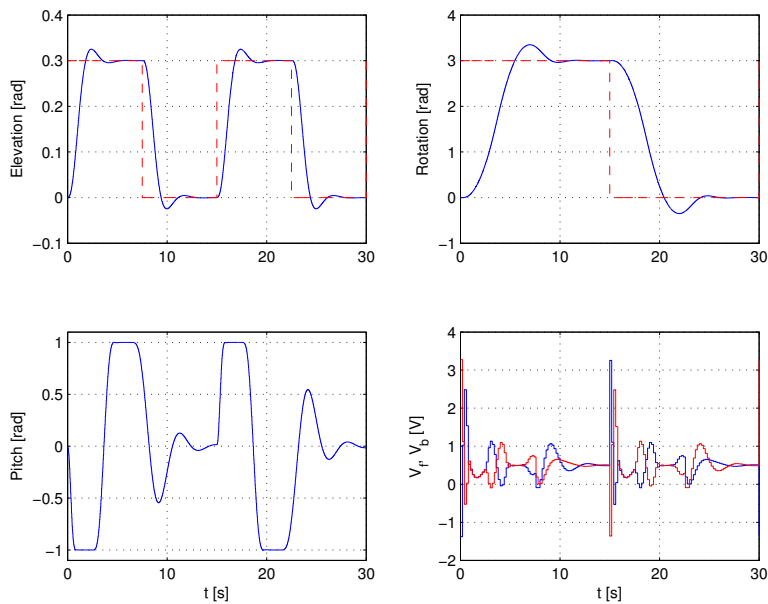


Figure 16.6: Response with  $N = 25$