

Order Full Rate, Leadtime Variability, and Advance Demand Information in an Assemble-To-Order System

by Lu, Song, and Yao (2002)

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This summary presentation is based on: Lu, Yingdong, and Jing-Sheng Song. "Order-Based Cost Optimization in Assemble-to-Order Systems." To appear in *Operations Research*, 2003.

Preview

- Assembly-to-order system
 - Each product is assembled from a set of components,
 - Demand for products following batch Poisson processes,
 - Inventory of each component follows a base-stock policy
 - Replenishment leadtime *i.i.d.* random variables for each component.
- Model as a $M^x / G / \infty$ queue, driven by a common multiclass batch Poisson input stream
 - Derive the joint queue-length distribution,
 - Order fulfillment performance measure.

Model

- M different components, and $F = \{1, 2, \dots, m\}$ are the component indices.
- Customer orders arrive as a stationary Poisson process, $\{A(t), t \geq 0\}$, with rate λ .
- Order type K : it contains positive units of component in K and 0 units in $F \setminus K$.
 - An order is of type K with probability q^K , $\sum_K q^K = 1$
 - Type K order stream forms a compound Poisson process with rate $\lambda^K = q^K \lambda$
 - A type- K order has Q_j^K units for each component j , $Q_K = (Q_j^K, j \in K)$ has a known discrete distribution.
- For each component i , the demand process forms a compound Poisson process.

Model

- Demand are filled on a FCFS basis.
- Demand are backlogged (if one or more components are missing), and are filled on a FCFS basis.
- Inventory of each component is controlled by an independent base-stock policy,
 - Where s_i is the base-stock level for component i
 - For each component i , replenishment leadtimes, L_i , are *i.i.d.*, with a cdf of G_i
 - Net inventory at time t , $I_i(t) = s_i - X_i(t)$, $i = 1, \dots, m$, where $X_i(t)$ is the number of outstanding orders of component i at time t .
- Immediate availability of all components needed for an arriving demand as the “off-the-shelf” fill rate.
 - Off-the-shelf fill rate of component i , $f_i = P[X_i + Q_i \leq s_i]$
 - Off-the-shelf fill rate of demand type K , $f^K = P[X_i + Q_i^K \leq s_i, \forall i \in K]$
 - Average (over all demand types) off-the-shelf fill rate, $f = \sum_K q^K f^K$

Performance Analysis

- Derive the joint distribution and steady state limit of vector $X(t) = (X_1(t), \dots, X_m(t))$

(See “Suppliers/Arrivals Replenishment Orders”
diagram in Lu, Song, and Yao paper)

- Each component i , the number of outstanding orders is exactly the number of jobs in service in an $M_i^{Q_i} / G_i / \infty$ queue with Poisson arrival λ_i and batch size Q_i
- The m queues are not independent.
- Given the number of demand arrivals up to t , the $X_i(t)$ s are independent of one another.

Performance Analysis

- Proposition 1: $X(t) = (X_1(t), \dots, X_m(t))$ has a limiting distribution. Derive the generating function of X .
- In the special case of unit arrival, $Q_i \equiv 1$, the generating function of X corresponds to a multivariate Poisson distribution. For each i , X_i is a Poisson variable with parameter $\lambda_i \ell_i = (\sum_{K \in \mathcal{R}_i} \lambda^K) \ell_i$
- The correlation of the queue is solely induced by the common arrivals. If the proportion of the demand types that require both i and j are very small, the correlation between X_i and X_j is negligible.
- Level of correlation is independent of the demand rate.
- Reducing the variability of leadtime or batch sizes will result in a higher correlation among the queue lengths of outstanding jobs.

Response-time-based order fill rate

- 1) $f^K(w)$ is the probability of having all the components ready within w units of time.
- 2) $D_i(t, t+u] := D_i(t+u) - D_i(t)$
- 3) Total number of departures from queue i in $(\tau, \tau+w) = X_i(\tau) + D_i(\tau, \tau+w) - X_i(\tau+w)$
- 4) $I_i(\tau) + \{X_i(\tau) + D_i(\tau, \tau+w) - X_i(\tau+w)\} \geq 0$
- 5) $X_i(\tau+w) - D_i(\tau, \tau+w) \leq s_i, i \in K$

- 6)
$$X_i(\tau+w) = X_i^w(\tau) + \sum_{n=1}^{Q_i^K} 1\{L_i^n > w\} + X_i(\tau, \tau+w)$$

- 7) Demand at τ can be supplied by $\tau+w$ iff

$$X_i^w(\tau) + X_i(\tau, \tau+w) - D_i(\tau, \tau+w) \leq s_i - \sum_{n=1}^{Q_i^K} 1\{L_i^n > w\}, i \in K$$

$$Y_i := X_i^w - Y_i^w$$

- 8) Order fill rate of type-K demand within time window w , $f^K(w) = P \left[Y_i + \sum_{n=1}^{Q_i^K} 1\{L_i^n > w\} \leq s_i, \forall i \in K \right]$

- 9) Mean:
$$E[Y_i] = \left(\sum_{\mathfrak{S} \in \mathfrak{R}_i} \lambda^{\mathfrak{S}} E(Q_i^{\mathfrak{S}}) \right) (\ell_i - w)$$

Connection to advance demand information

- Suppose each order arrival epoch is known w time units in advance, where $w > 0$ is a deterministic constant.
- Suppose a type- K order arrives at τ , and this information is known at $\tau + w$, we can fill this order upon its arrival with probability,

$$f_A^K(0) = \mathbb{P} \left\{ X_i^w(\tau - w) + X_i(\tau - w, \tau] - D_i(\tau - w, \tau] \leq s_i - \sum_{n=1}^{Q_i^K} \mathbf{1}[L_i^n > w], i \in K \right\} = f^K(w).$$

- Advance demand information improves the off-the-shelf fill rate: $f^K(w) \geq f^K(0)$
 - Compare $f_A^K(0)$ with that of the modified system, $\hat{f}^K(0)$, where leadtime is reduced from L_i to $\hat{L}_i = [L_i - w]^+$
- $$\hat{f}^K(0) = \mathbb{P} \left\{ X_i^w + \sum_{n=0}^{Q_i^K} \mathbf{1}[\hat{L}_i^n > 0] \leq s_i, \forall i \in K \right\} \leq f^K(w)$$
- Knowing demand in advance (by w time units) is more effective, in terms of order fill rate, than reducing the supply leadtime of components.